

Adjunctions, limits

&

colimits

Def An adjunction  $A \perp B$  between  $\infty$ -categories is comprised of:

- a pair of  $\infty^u$ -categories  $A$  &  $B$
- a pair of  $\infty$ -functors

$$u: A \rightarrow B \quad \& \quad f: B \rightarrow A$$

- a pair of  $\infty$ -natural transformations

$$\eta: id_B \Rightarrow uf \quad \& \quad \varepsilon: fu \Rightarrow id_A$$

So that

$$\begin{array}{c}
 B \xlongequal{\quad} B \\
 \begin{array}{ccc}
 \nearrow u & & \searrow f \\
 A & \xlongequal{\quad} & A \\
 \varepsilon \Downarrow & & \Downarrow \eta \\
 \end{array}
 \end{array}
 =
 \begin{array}{c}
 \begin{array}{ccc}
 & \nearrow B & \nwarrow \\
 u & (=) & u \\
 & \searrow A & \nearrow
 \end{array}
 \\
 u \xrightarrow{\eta^-} ufu \xrightarrow{\varepsilon^-} u
 \end{array}$$

$$\begin{array}{c}
 B \xlongequal{\quad} B \\
 \begin{array}{ccc}
 \searrow f & & \nearrow f \\
 \Downarrow \eta & & \Downarrow \varepsilon \\
 A & \xlongequal{\quad} & A
 \end{array}
 \end{array}
 =
 \begin{array}{c}
 \begin{array}{ccc}
 & \nwarrow B & \nearrow \\
 f & (=) & f \\
 & \searrow A & \nearrow
 \end{array}
 \\
 f \xrightarrow{\eta^-} fuf \xrightarrow{\varepsilon^-} f
 \end{array}$$

Remark. In the the setting of  $(\infty, n)$ - or  $(\infty, \infty)$ -categories  
this is "pseudo-style" adjunction

↑  
It is not most general adjunctions

But: its relationships to the equivalences  
&  
to the notions of (co-)limits

Lemma. An adjunction in a 2-category is preserved by any 2-functor

Example. Adjunction between 1-cats  
 $\text{Cat} \hookrightarrow \text{hQCats}$

$$A \overset{\perp}{\rightleftarrows} B \rightsquigarrow \mathcal{N}(A) \overset{\perp}{\rightleftarrows} \mathcal{N}(B)$$

↑  
regarded as a nerve

Example. Quillen adjunctions

Prop. Given an adjunction  $A \begin{array}{c} \xleftarrow{f} \\ \perp \\ \xrightarrow{u} \end{array} B$  between  $\infty$ -cats

Then

•  $\forall \infty$ -cat  $X$

$$\text{Fun}(X, A) \begin{array}{c} \xleftarrow{f_*} \\ \perp \\ \xrightarrow{u_*} \end{array} \text{Fun}(X, B)$$

•  $\forall \infty$ -cat  $X$

$$\text{hFun}(X, A) \begin{array}{c} \xleftarrow{f_*} \\ \perp \\ \xrightarrow{u_*} \end{array} \text{hFun}(X, B)$$

•  $\forall$  simplicial set  $\mathcal{U}$

$$A^{\mathcal{U}} \begin{array}{c} \xleftarrow{f^{\mathcal{U}}} \\ \perp \\ \xrightarrow{u^{\mathcal{U}}} \end{array} B^{\mathcal{U}}$$

• If the ambient  $\infty$ -cosmos is cartesian closed, then  $\forall \infty$ -cat  $\mathcal{C}$

$$A^{\mathcal{C}} \begin{array}{c} \xleftarrow{\quad} \\ \perp \\ \xrightarrow{\quad} \end{array} B^{\mathcal{C}}$$

Prop. Adjunctions compose

$$\begin{array}{ccc} \begin{array}{ccc} C & \xrightarrow{f'} & B \\ \downarrow \perp & & \downarrow \perp \\ C & \xrightarrow{f} & A \\ \uparrow u' & & \uparrow u \end{array} & \rightsquigarrow & \begin{array}{ccc} C & \xrightarrow{ff'} & A \\ \downarrow \perp & & \downarrow \perp \\ C & \xrightarrow{f} & A \\ \uparrow u'u & & \uparrow u \end{array} \end{array}$$

Prop. (uniqueness of adjoints)

- If  $f \dashv u$  &  $f' \dashv u \Rightarrow f \cong f'$
- Conversely, if  $f \dashv u$  and  $f \cong f' \Rightarrow f' \dashv u$

Lemma. (minimal adjunction data)

$$\begin{array}{c} \xrightarrow{f} \\ A \quad \perp \quad B \\ \xleftarrow{u} \end{array} \iff \exists \text{ nat. transf. } \begin{array}{l} \text{id}_B \Rightarrow uf \\ fu \Rightarrow \text{id}_A \end{array} \text{ so that}$$

the triangle equality composites are invertible:

$$f \Rightarrow fuf \Rightarrow f \quad \& \quad u \Rightarrow ufu \Rightarrow u$$

Proof:  $(\Rightarrow)$  Obvious

$$(\Leftarrow) \cdot \eta: \text{id}_B \Rightarrow uf \quad \& \quad \varepsilon': fu \Rightarrow \text{id}_A$$

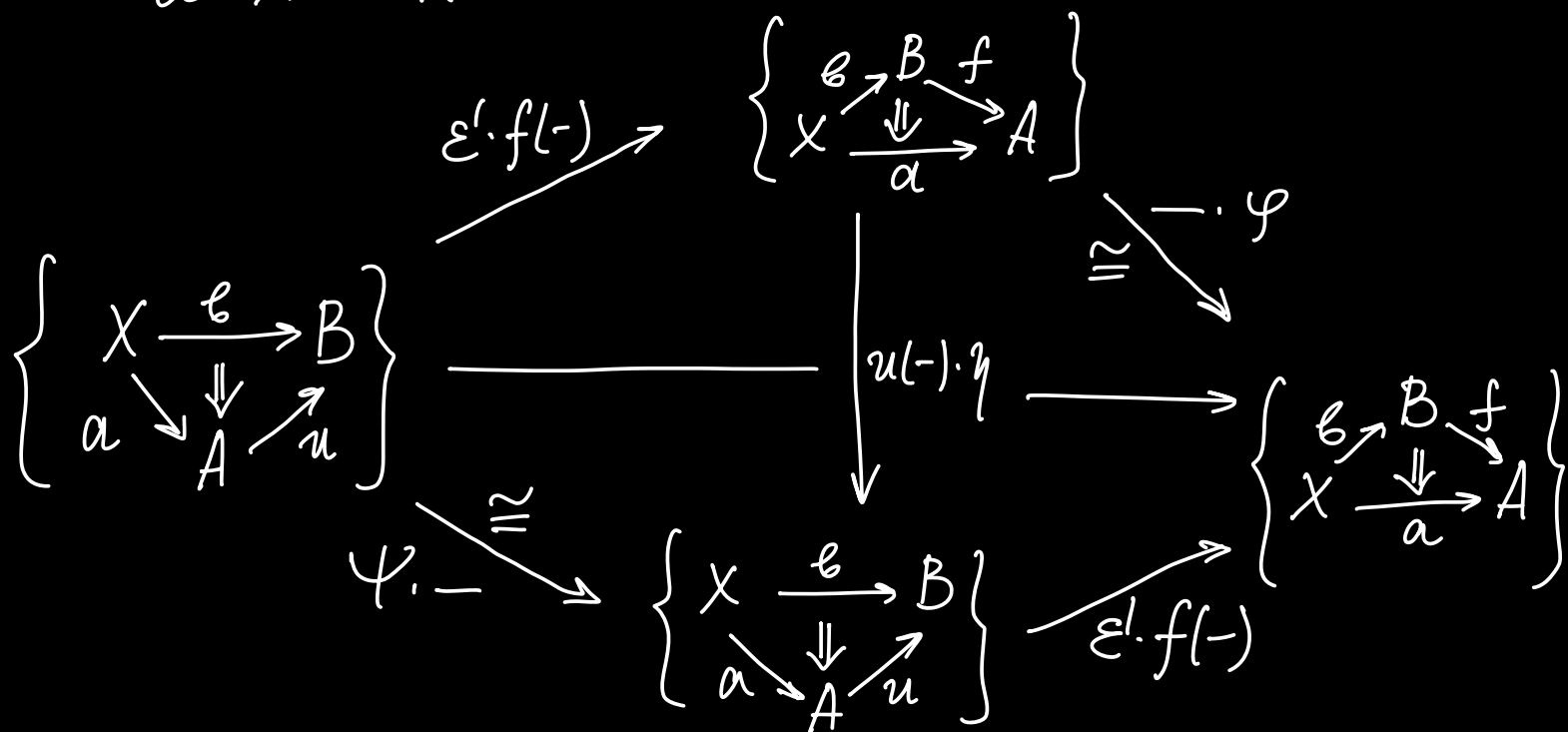
$$\varphi := f \xrightarrow{f\eta} fuf \xrightarrow{\varepsilon'f} f \quad \& \quad \psi := u \xrightarrow{\eta u} ufu \xrightarrow{u\varepsilon'} u$$

are isomorphisms

• Construct an adjunction  $f \dashv u$  with unit  $\eta$  and counit  $\varepsilon'$

$$\bullet \quad b \Rightarrow ua \xrightarrow{f \circ -} fb \Rightarrow fua \xrightarrow{\varepsilon'} fb \Rightarrow fua \xrightarrow{\varepsilon'} a \rightsquigarrow fb \Rightarrow a$$

• Let  $b: X \rightarrow B$   
 $a: X \rightarrow A$  be fixed generalized elements





• By 2-of-6 property, we have all six morphisms being  
bijections

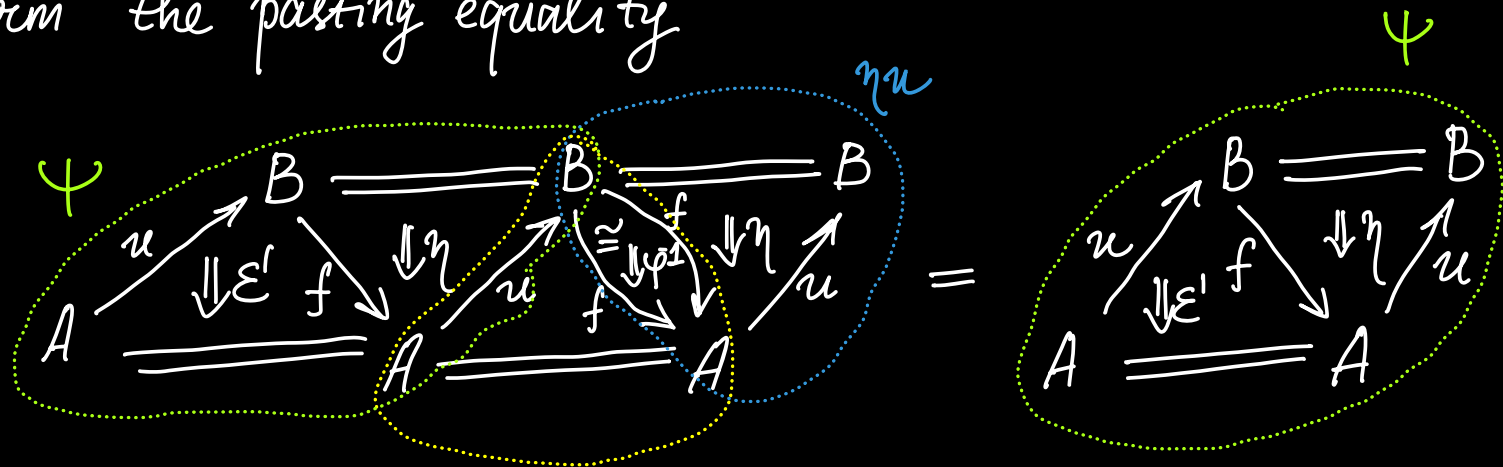
• Define

$$\varepsilon := \begin{array}{ccc} & B & \\ \nearrow u & & \searrow f \\ A & \xrightarrow{\varepsilon'} & A \end{array} \quad \begin{array}{c} \Downarrow \eta \\ \Downarrow \varepsilon' \\ \Downarrow \psi^{-1} \end{array}$$

so that

$$\begin{array}{ccc} B & \xrightarrow{\quad} & B \\ f \searrow & \Downarrow \eta & \nearrow u \\ & A & \searrow f \\ & \xrightarrow{\quad} & A \end{array} \quad \begin{array}{c} \Downarrow \varepsilon' \\ \Downarrow \varepsilon \end{array} = \begin{array}{ccc} B & \xrightarrow{\quad} & B \\ \eta \searrow & \Downarrow \eta & \nearrow u \\ & A & \searrow f \\ & \xrightarrow{\quad} & A \end{array} \quad \begin{array}{c} \Downarrow \varepsilon' \\ \Downarrow \varepsilon \\ \Downarrow \psi^{-1} \end{array} = f \left( \begin{array}{c} B \\ = \\ A \end{array} \right) f$$

• Form the pasting equality



$$(u\epsilon \cdot \eta u) \cdot \Psi = \Psi$$

$$\Psi := u \xRightarrow{\eta u} u f u \xRightarrow{u \epsilon'} u$$

• But  $\Psi$  is invertible

$$\text{So, } u\epsilon \cdot \eta u = \text{id}_u$$



Corollary. (adjoint equivalences) Any equivalence can be promoted to an adjoint equivalence by modifying one of the 2-cells, i.e., the invertible 2-cells in an equivalence can be chosen so as to satisfy the triangle equalities

$$f, g \text{ are inverse equiv.} \rightsquigarrow f \dashv g \ \& \ g \dashv f$$

Proof:

$$\begin{array}{ccc}
 \begin{array}{ccc}
 & f & \\
 A & \xrightarrow{\quad} & B \\
 & \perp & \\
 & g & \\
 & \xleftarrow{\quad} & 
 \end{array}
 & 
 \begin{array}{ccc}
 & \cong & \\
 A & \xrightarrow{\quad} & A \\
 & \Downarrow \alpha & \\
 & gf & 
 \end{array}
 & 
 \begin{array}{ccc}
 & fg & \\
 B & \xrightarrow{\quad} & B \\
 & \cong & \\
 & \Downarrow \beta & 
 \end{array}
 \end{array}$$

$$f \rightsquigarrow fgf \rightsquigarrow f \quad \& \quad g \rightsquigarrow gfg \rightsquigarrow g$$

So, by the previous lemma  $f$  and  $g$  fit into an adjunction  $\triangleleft$

Prop. Adjunctions are equivalence invariant

$$u: A \rightarrow B \text{ admits a left adjoint} \iff \forall \begin{matrix} A' \cong A \\ B' \cong B \end{matrix}$$

$\infty$ -cats

the functor

$$u': A' \rightarrow B'$$

admits a left adjoint

Proof:  $\exists f \dashv u$  with  $A' \cong A, B' \cong B$  - adjoint equiv.

$$\begin{array}{ccccc} A' & \xleftarrow{\sim} & A & \xleftarrow{f} & B & \xleftarrow{\sim} & B' \\ \xrightarrow{\sim} & \perp & \xrightarrow{u} & \perp & \xrightarrow{\sim} & \perp & \xrightarrow{\sim} \\ & & & & & & \end{array}$$

Conversely, if  $u': A' \xrightarrow{\sim} A \xrightarrow{u} B \xrightarrow{\sim} B'$  admits a left adj  $f'$  then

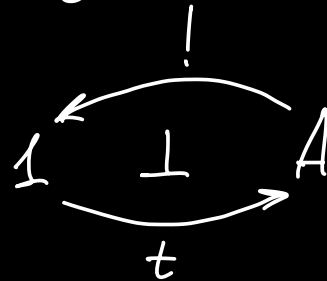
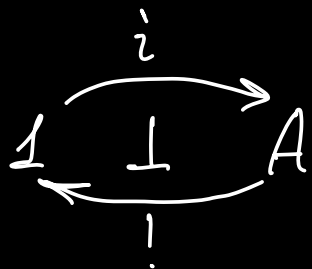
$$f \Rightarrow fuf \Rightarrow f \quad \& \quad u \Rightarrow ufu \Rightarrow u$$

$$\begin{array}{ccccccc} & & \xleftarrow{\sim} & \xleftarrow{f'} & \xleftarrow{\sim} & \xleftarrow{\sim} & \xleftarrow{\sim} \\ & & A & A' & A & B & B' \\ \xrightarrow{\sim} & \perp & \xrightarrow{u} & \xrightarrow{\sim} & \xrightarrow{u} & \xrightarrow{\sim} & \xrightarrow{\sim} \\ & & & & & & \end{array}$$



# Initial & terminal elements

Def. An initial element in an  $\infty$ -cat  $A$  is a left adjoint to the unique functor  $! : A \rightarrow \mathbb{1}$   
 A terminal element is a right adj.



Lemma. <sup>(minimal data)</sup> To define an initial element in an  $\infty$ -cat  $A$ , it suffices to specify:

- an element  $i : \mathbb{1} \rightarrow A$

- a nat. transf.  $A \begin{array}{c} \xrightarrow{!} \mathbb{1} \xrightarrow{i} A \\ \xrightarrow{=} \mathbb{1} \xrightarrow{i} A \\ \Downarrow \varepsilon \end{array} A$  so that  $\varepsilon_i : i \Rightarrow i$  in  $hA$  is unim.

Proof: •  $\infty$ -cat  $\mathcal{K} \in \mathcal{K}$  is 2-terminal in the homotopy 2-cat  $h\mathcal{K}$

•  $i: \mathbb{1} \rightarrow A$  is a section of  $!: A \rightarrow \mathbb{1}$

So,  $!i = id_{\mathbb{1}}$  &  $\eta: \mathbb{1} \xrightarrow{\sim} !i$  is an iso

We have a triangle

$$\begin{array}{c}
 A \xrightarrow{\quad} A \\
 \begin{array}{ccc}
 \nearrow i & & \searrow ! \\
 \eta \Downarrow & & \Downarrow \varepsilon \\
 \mathbb{1} & \xrightarrow{\quad} & \mathbb{1}
 \end{array}
 \end{array}
 \begin{array}{c}
 \nearrow i \\
 (=) \\
 \searrow i \\
 \mathbb{1}
 \end{array}$$

$$i \xrightarrow{i\eta} i!i \xrightarrow{\varepsilon i} i$$

$$\varepsilon: i! \rightarrow id$$

$$\eta: id \xrightarrow{\sim} !i$$

this does not

We want:

$$\varepsilon i = id_i$$

$$\begin{array}{c}
 A \xrightarrow{\quad} A \\
 \begin{array}{ccc}
 \searrow i & & \nearrow ! \\
 \Downarrow \varepsilon & & \Downarrow \eta \\
 \mathbb{1} & \xrightarrow{\quad} & \mathbb{1}
 \end{array}
 \end{array}
 \begin{array}{c}
 \searrow i \\
 (=) \\
 \nearrow ! \\
 \mathbb{1}
 \end{array}$$

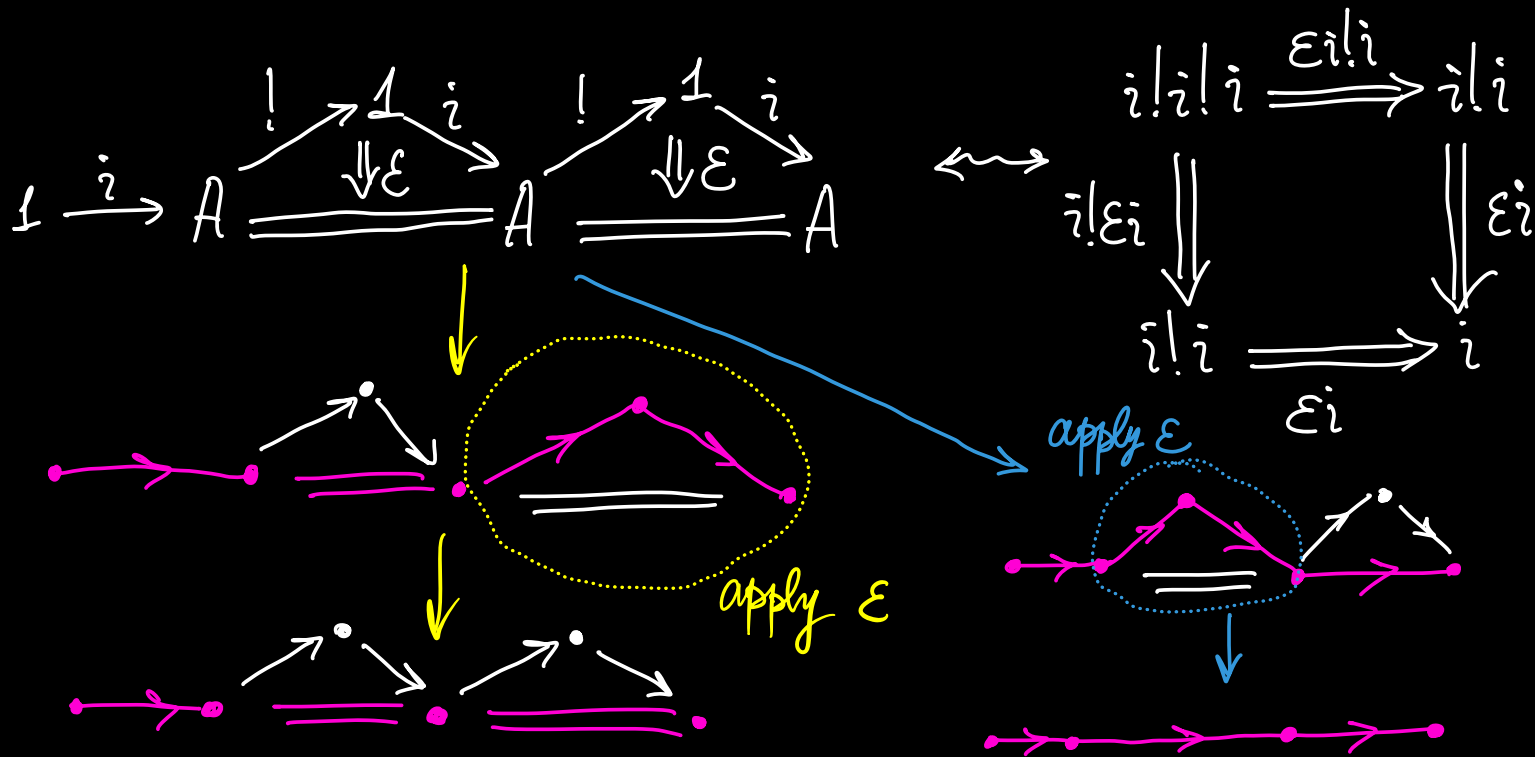
$$!i \xrightarrow{\eta!} !i! \xrightarrow{! \varepsilon} !$$

this holds automatically

$$! \varepsilon(a) = !(a)$$

- To do this, it suffices to require an isomorphism
 
$$\varepsilon_i : \bar{i} \cong i \quad (\text{See the prop. above})$$

- Consider the horizontal composite and represent it as a vertical composite in two ways:



• So,  $\varepsilon_i \cdot \varepsilon_i = \varepsilon_i \implies \varepsilon_i = \text{id}_i$  since  $\varepsilon_i$  is an isomorphism  
by the assumption  $\triangleleft$

Lemma (uniqueness) Any two initial elements in an  $\infty$ -cat  $A$  are isomorphic in  $hA$  and if  $a \cong i$  in  $hA \implies a$  is initial too  $\triangleleft$

Proof:  $\forall$  left adjoints  $i$  and  $i'$  to  $! : A \rightarrow \mathbb{1}$   
are nat. isomorphic

•  $\forall a : \mathbb{1} \rightarrow A$  that is isomorphic to left adj to  $! : A \rightarrow \mathbb{1}$   
is itself a left adjoint

• A functor isomorphism  $i \cong i'$  gives  $i \cong i'$  in  $hA$   $\triangleleft$



Lemma. An element  $i: 1 \rightarrow A$  is initial  $\Leftrightarrow \forall f: X \rightarrow A$   
 $\exists$  a unique 2-cell with boundary

$$\begin{array}{ccc}
 & 1 & \\
 X \xrightarrow{i!} & \downarrow \exists! & \xrightarrow{i} A \\
 & f & 
 \end{array}$$

Proof. •  $i \dashv ! \rightsquigarrow \mathbb{1} \cong \mathbf{hFun}(X, \mathbb{1}) \begin{array}{c} \xrightarrow{i_*} \\ \perp \\ \xleftarrow{\quad} \end{array} \mathbf{hFun}(X, A)$

So, the constant functor  $i!: X \rightarrow A$  is initial in  $\mathbf{hFun}(X, A)$

• Conversely, if  $i: 1 \rightarrow A$  satisfies univ. property of the statement

Then apply it to the  $\text{id}_A: A \rightarrow A$  any by lemma above

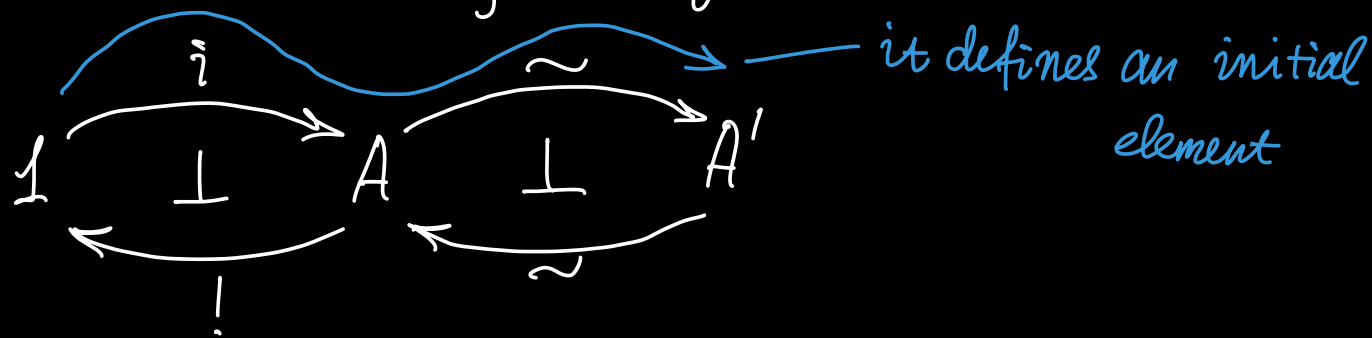
we are done ◻

$$\begin{array}{ccc}
 & i & \\
 & \curvearrowright & \\
 \mathbb{1} & \xleftarrow{\quad} & hA \\
 & \curvearrowleft & \\
 & t & 
 \end{array}$$

- Remark.
- Being homotopy initial is weaker than being initial in the  $\infty$ -cat
  - But a homotopy initial element in a complete  $(\infty, 1)$ -cat defines an initial element

Lemma. If  $A$  has an initial element and  $A \cong A'$  then  $A'$  has an initial element and they are respected by the equivalence up to isomorphism

Proof:  $A \cong A' \rightsquigarrow$  an adjoint equivalence



By the uniqueness of initial elements, the equivalence  $A' \xrightarrow{\sim} A$  preserves initial elements △

## Limits & colimits

(co-)limits should be interpreted as homotopy ones in  $\infty$ -cat

(Co-)limits of a diagram valued inside an  $\infty$ -cat  $A$  in some  $\infty$ -cosmos

- indexed by a simpl. set  $J$  in an  $\infty$ -cat  $A$  in a generic  $\infty$ -cosmos
- indexed by an  $\infty$ -cat  $J$  and valued in an  $\infty$ -cat  $A$  in a cartesian closed  $\infty$ -cosmos

Def. (diagram  $\infty$ -cat)

For  $A \in \mathcal{K}$ ,  $J \in \mathbf{sSet}$

or

$J \in \mathcal{K}_{\text{cart. closed}}$

define a diagram of shape

$J$  in  $A$  as an element

$$d: \mathbb{1} \rightarrow A^J$$

$$\mathcal{S}\text{Set}^{\mathcal{J}} \times \mathcal{K} \longrightarrow \mathcal{K}$$

$$(\mathcal{J}, A) \longmapsto A^{\mathcal{J}}$$

$$\mathcal{K}^{\text{op}} \times \mathcal{K} \longrightarrow \mathcal{K}$$

$$(\mathcal{J}, A) \longmapsto A^{\mathcal{J}}$$

— bifunctors

Apply bifunctor to  $! : \mathcal{J} \rightarrow 1$

$$\Delta : A \longrightarrow A^{\mathcal{J}} \text{ — constant diagram functor}$$

Def. (limit & colimit functor)

$$\begin{array}{ccc}
 & \text{colim} & \\
 A & \xrightarrow{\quad} & A \\
 & \Delta & \\
 & \perp & \\
 & \text{lim} & \\
 & \xleftarrow{\quad} & 
 \end{array}$$

Lemma. Products or coproducts in an  $\infty$ -cat  $A$  also define ones in its homotopy cat  $hA$

Proof: . If  $J$  is a set  $\Rightarrow A^J \cong \prod_J A$

$$\begin{array}{ccc} h\mathcal{K} & \xrightarrow{h\text{Fun}(1, -)} & \text{Cat} \\ A & \longmapsto & hA \end{array} \quad \leftarrow \text{preserves products}$$

$$h(A^J) \cong h\left(\prod_J A\right) \cong \prod_J hA \cong (hA)^J \quad \begin{array}{l} J \text{ is a set} \end{array}$$

So,

$$(hA)^J \cong h(A^J) \begin{array}{c} \xrightarrow{\text{colim}} \\ \perp \\ \xrightarrow{\Delta} \\ \perp \\ \xrightarrow{\text{lim}} \end{array} hA \quad \text{Put } J = \emptyset \triangleleft$$

The definition is insufficiently general!

Def. (absolute lifting diagram)

Given a cospan  $C \xrightarrow{g} A \xleftarrow{f} B$  in a 2-cat

An absolute left lifting of  $g$  through  $f$  is given by

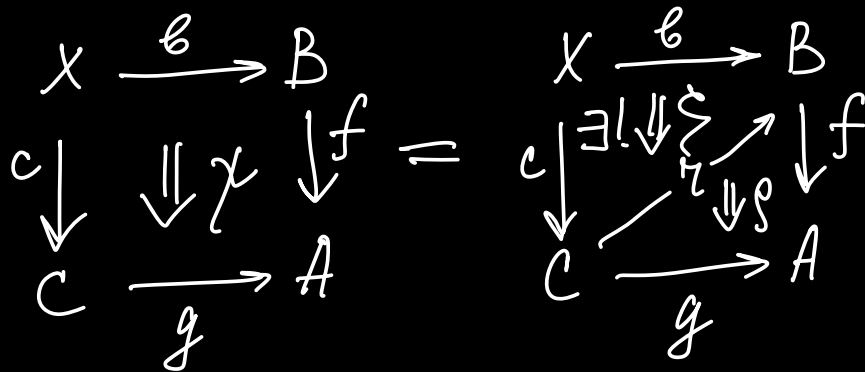
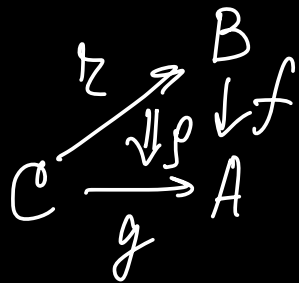
a 1-cell  $l$  & 2-cell  $\lambda$ :

$$\begin{array}{ccc}
 & & B \\
 & \nearrow l & \downarrow f \\
 C & \xrightarrow{g} & A
 \end{array}
 \quad \text{s.t.}$$

$$\begin{array}{ccc}
 X & \xrightarrow{e} & B \\
 c \downarrow & \Uparrow \chi & \downarrow f \\
 C & \xrightarrow{g} & A
 \end{array}
 =
 \begin{array}{ccc}
 X & \xrightarrow{\quad} & B \\
 \downarrow \exists! \Uparrow \xi & \nearrow l & \downarrow f \\
 C & \xrightarrow{\quad} & A
 \end{array}$$

Any 2-cell  $\chi$   
factors through  
( $l, \lambda$ )

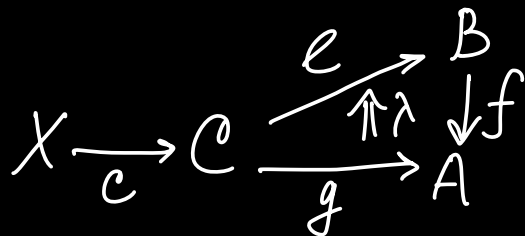
An absolute right lifting of  $g$  through  $f$ :



"Absolute" means

Lemma. Left (right) lifting diagrams are stable under restriction of their domain object:

If  $(\ell, \lambda)$  - abs. left lifting of  $g$  through  $f$ , then  $\forall c: X \rightarrow C$   
 $(\ell c, \lambda c)$  defines an absolute left lifting of  $gc$  through  $f$

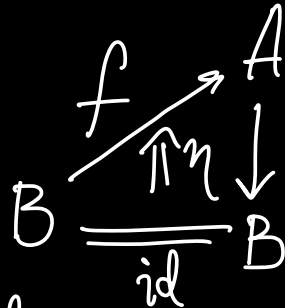




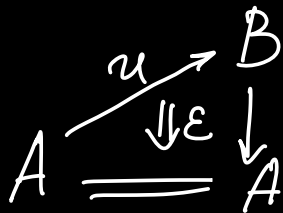
Example:  $id_B \Rightarrow uf$  is the unit of  $f \dashv u$



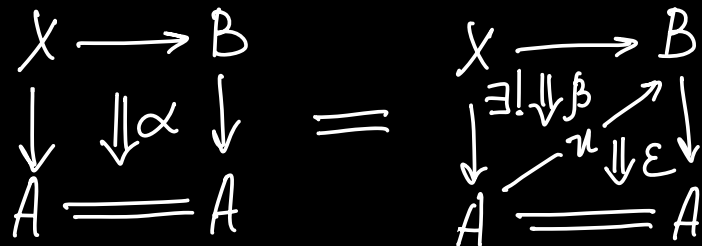
$(f, \eta)$  defines an absolute left lifting diagram



Dually for  $\varepsilon: fu \Rightarrow id_A$



Proof:



$$\alpha: fb \Rightarrow a \rightsquigarrow \beta: b \Rightarrow ua$$

$$\begin{array}{ccc}
 \beta & \xleftarrow{\quad} & \alpha \\
 \uparrow & & \uparrow \\
 \text{hFun}(X, B) & \begin{array}{c} \xrightarrow{f_*} \\ \perp \\ \xleftarrow{u_*} \end{array} & \text{hFun}(X, A)
 \end{array}$$

If  $f \dashv u$  then it is an adjunction and  $(u, \epsilon)$  defines an absolute right lifting of  $\text{id}_A$  through  $f$

Conversely, the unit & triangle identities of an adj. can be extracted from the univ. prop. of the right lifting diagram  $\triangleleft$

In particular,  $\text{colim} \dashv \Delta \dashv \text{lim}$  define

$$\begin{array}{ccc}
 \begin{array}{ccc}
 & \text{colim} \rightarrow & A \\
 A^{\mathcal{J}} & \xrightarrow{\eta} & \downarrow \Delta \\
 A & \xlongequal{\quad} & A^{\mathcal{J}}
 \end{array} & \& & \begin{array}{ccc}
 & \rightarrow & A \\
 A^{\mathcal{J}} & \xrightarrow{\epsilon} & \downarrow \\
 A & \xlongequal{\quad} & A^{\mathcal{J}}
 \end{array}
 \end{array}$$

These univ. properties are preserved under restriction. It motivates

Def. (limit & colimit) A colimit of a family of diagrams  $d: \mathcal{D} \rightarrow A^{\mathcal{J}}$  of shape  $\mathcal{J}$  in an  $\infty$ -cat  $A$  is given by an absolute left lifting diagram

$$\begin{array}{ccc}
 & \text{colim } d & \rightarrow A \\
 \mathcal{D} & \xrightarrow{\uparrow \eta} & \downarrow \Delta \\
 & \rightarrow & A^{\mathcal{J}}
 \end{array}$$

comprised of a generalized element  $\text{colim } d: \mathcal{D} \rightarrow A$  & a colimit cone  $\eta: d \Rightarrow \Delta \text{colim } d$

Dually,

$$\begin{array}{ccc}
 & \text{lim } d & \rightarrow A \\
 \mathcal{D} & \xrightarrow{\downarrow \varepsilon} & \downarrow \Delta \\
 & \xrightarrow{d} & A^{\mathcal{J}}
 \end{array}$$

$\varepsilon: \Delta \text{lim } d \Rightarrow d$  — a limit cone

a limit of a family of diagrams  $d: \mathcal{D} \rightarrow A^{\mathcal{J}}$  of shape  $\mathcal{J}$  in an  $\infty$ -cat  $A$

Example An initial element  $i: 1 \rightarrow A$  can be regarded as a colimit of the empty diagram

The constant diagram functor:

$$! : A \longrightarrow A^{\mathbb{J}} = A^{\emptyset} = 1$$

$$\begin{array}{ccc}
 \begin{array}{ccc} & & A \\ & \nearrow i & \downarrow ! \\ 1 & \xlongequal{\quad} & 1 \end{array} & 
 \begin{array}{ccc} X & \xrightarrow{f} & A \\ \downarrow ! & \uparrow \gamma & \downarrow ! \\ 1 & \xlongequal{\quad} & 1 \end{array} & = & 
 \begin{array}{ccc} X & \xrightarrow{f} & A \\ \downarrow ! & \uparrow \exists! \Sigma & \downarrow ! \\ 1 & \nearrow i & \downarrow ! \\ & \xlongequal{\quad} & 1 \end{array}
 \end{array}$$

We want: an initial element defines an absolute left lifting diagram whose 2-cell is identity

But the existence & uniqueness of  $\Sigma$  follow from initiality of  $i: 1 \rightarrow A$  among all generalized elements  $f: X \rightarrow A$

Example (exercise) In a cartesian closed  $\infty$ -category,  $i: 1 \rightarrow A$  can be regarded as a limit of  $\text{id}_A: A \rightarrow A$

Theorem Right adjoints preserve limits

Proof:  $\square$   $A$  admits limits of  $d: 1 \rightarrow A^J$

$$\begin{array}{ccc}
 & \lim d & \rightarrow A \\
 1 & \nearrow & \downarrow \Delta \\
 & d & \rightarrow A^J
 \end{array}$$

•  $f \dashv u \rightsquigarrow f^J \dashv u^J$

• Show that

$$\begin{array}{ccccc}
 \lim d & \rightarrow & A & \xrightarrow{u} & B \\
 1 & \nearrow & \downarrow \Delta & & \downarrow \Delta \\
 & d & \rightarrow & A^J & \rightarrow & B^J
 \end{array}$$

it is again an absolute right lifting diagram

• Consider

$$\begin{array}{ccccc}
 X & \xrightarrow{\ell} & B & & \\
 \downarrow ! & & \Downarrow \gamma & & \downarrow \Delta \\
 1 & \xrightarrow{d} & A^J & \xrightarrow{\kappa^J} & B^J
 \end{array}$$

• Add a square to the right hand side, comparing with  $f$

$$\begin{array}{ccccccc}
 X & \xrightarrow{\ell} & B & \xrightarrow{f} & A & & \\
 \downarrow ! & & \Downarrow \gamma & & \downarrow \Delta & & \\
 1 & \xrightarrow{d} & A^J & \xrightarrow{\kappa^J} & B^J & \xrightarrow{f^J} & A^J \\
 & & & & \Downarrow \varepsilon^J & & \\
 & & & & A^J & & 
 \end{array}
 \quad = \quad
 \begin{array}{ccccccc}
 X & \xrightarrow{\ell} & B & \xrightarrow{f} & A & & \\
 \downarrow ! & & \Downarrow \xi & & \downarrow \Delta & & \\
 1 & \xrightarrow{d} & A^J & \xrightarrow{\text{land}} & A^J & & \\
 & & & & \downarrow \lambda & & \\
 & & & & A^J & & 
 \end{array}$$

just compose & factorize

- Add a square to the right hand side, comparing with  $u$

$$\begin{array}{ccc}
 X \xrightarrow{e} B \xrightarrow{f} A \xrightarrow{u} B & & X \xrightarrow{e} B \xrightarrow{f} A \xrightarrow{u} B \\
 \downarrow \exists! \downarrow \zeta & \Downarrow \eta & \downarrow \Delta \\
 1 \xrightarrow{d} A^J \xrightarrow{u^J} B^J & & 1 \xrightarrow{d} A^J \xrightarrow{u^J} B^J \\
 \downarrow \Delta & & \downarrow \Delta \\
 A^J \xrightarrow{u^J} B^J & & A^J \xrightarrow{u^J} B^J \\
 \downarrow \Delta & & \downarrow \Delta \\
 B^J & & B^J
 \end{array}$$

$$\begin{array}{ccc}
 X \xrightarrow{e} B \xrightarrow{=} B & & X \xrightarrow{e} B \\
 \downarrow \Delta & & \downarrow \Delta \\
 1 \xrightarrow{d} A^J \xrightarrow{u^J} B^J \xrightarrow{\eta^J} A^J \xrightarrow{u^J} B^J & & 1 \xrightarrow{d} A^J \xrightarrow{u^J} B^J \\
 \downarrow \Delta & & \downarrow \Delta \\
 A^J \xrightarrow{u^J} B^J & & A^J \xrightarrow{u^J} B^J \\
 \downarrow \Delta & & \downarrow \Delta \\
 B^J & & B^J
 \end{array}$$

apply a triangle identity of  $f^J \dashv u^J$

So,  $\gamma$  factors through the composite of  $\zeta$  and  $\eta$

Uniqueness is left  $\triangleleft$  as an exercise

Corollary Equivalences preserve limits

Moreover,

Prop. An equivalence  $f: A \xrightarrow{\sim} B$  preserves, reflects and creates limits and colimits

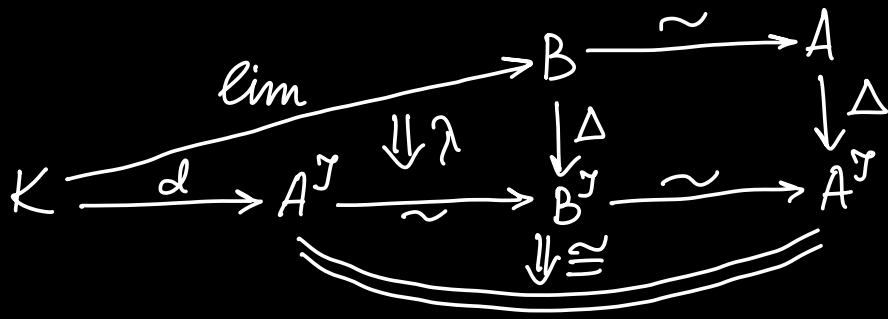
Proof: Prove that  $f$  reflects limits and colimits

- Consider a family of diagrams  $d: K \rightarrow A^J$  that admits limits in  $B$  after comp. with  $A \cong B$

$$\begin{array}{ccccc} & & & & B \\ & & \lim & \nearrow & \\ K & & & & \downarrow \Delta \\ & \xrightarrow{d} & A^J & \xrightarrow{\sim} & B^J \\ & & \Downarrow \lambda & & \end{array}$$

- The composite 2-cell





again defines  
an absolute  
right lifting  
diagram

The univ. prop. can be verified



Thank you!