

## The 2-category of quasi-categories

- The adjunction

$$h: \mathbf{sSet} \begin{array}{c} \xrightarrow{\quad} \\ \perp \\ \xleftarrow{\quad} \end{array} \mathbf{Cat} : \mathcal{N}$$

- The counit is an isomorphism

Lemma. The functor  $h: \mathbf{sSet} \rightarrow \mathbf{Cat}$  preserves finite product

▷ •  $(h-)$  &  $h(- \times -)$  preserve colimits

•  $\mathbf{sSet}$  &  $\mathbf{Cat}$  are cartesian closed

$$\bullet h\Delta^n \times h\Delta^m \stackrel{?}{\cong} h(\Delta^n \times \Delta^m)$$

•  $\Delta^n = \mathcal{N}(\tilde{n})$  where  $\tilde{n}$  is some category  $h\mathcal{N} = \mathcal{E}$  is iso

$$\bullet (h\Delta^n) \times (h\Delta^m) \cong (h\mathcal{N}\tilde{n}) \times (h\mathcal{N}\tilde{m}) \cong \tilde{n} \times \tilde{m} \cong h\mathcal{N}(\tilde{n} \times \tilde{m}) \cong h(\mathcal{N}\tilde{n} \times \mathcal{N}\tilde{m})$$

• Hence,  $\mathcal{h}$  and  $\mathcal{N}$  are strong monoidal

•  $h_* : \mathcal{S}\text{Cat} \begin{array}{c} \xrightarrow{\quad} \\ \perp \\ \xleftarrow{\quad} \end{array} \mathcal{Q}\text{Cat} : \mathcal{N}_*$

Def.  $q\text{Cat}_\infty \hookrightarrow \underline{\text{SSet}} \rightsquigarrow q\text{Cat}_2 := h_* q\text{Cat}_\infty$   
*2-category of quasi-categories*

$$\text{Ob}(q\text{Cat}_2) = \{\text{quasi-categories}\}$$

$$1\text{-cells of } q\text{Cat}_2 = \{\text{maps of quasi-categories}\}$$

$$2\text{-cells of } q\text{Cat}_2 = \{\text{homotopy classes of homotopies}\}$$

$$2\text{-cell } \alpha : f \Rightarrow g \rightsquigarrow 1\text{-simplex } \tilde{\alpha} : f \rightarrow g \text{ in } \mathcal{Y}^X$$

$$\alpha_1 \sim \alpha_2 \iff \tilde{\alpha}_1 \sim \tilde{\alpha}_2 \text{ as } 1\text{-simplices in } \mathcal{Y}^X$$

Prop.  $q\text{Cat}_2$  is cartesian closed

•  $h: q\text{Cat} \rightarrow \text{Cat} \rightsquigarrow 2\text{-functor } h_{\downarrow}: q\text{Cat}_2 \rightarrow \underline{\text{Cat}}$

$$\begin{array}{ccc} h(Y^X) \times hX \cong h(Y^X \times X) & \xrightarrow{h(\text{ev})} & hY \\ & \downarrow \text{adj} & \\ & h(Y^X) & \rightarrow hY^{hX} \end{array} \quad \left. \vphantom{\begin{array}{ccc} h(Y^X) \times hX \cong h(Y^X \times X) & \xrightarrow{h(\text{ev})} & hY \\ & \downarrow \text{adj} & \\ & h(Y^X) & \rightarrow hY^{hX} \end{array}} \right\} \text{On hom-cats}$$

- The aim: the category theory of quasi-categories
- $q\text{Cat}_2$  has finite products (see above)
- 2-limits theory
- Cotensors with the walking arrow category  $\mathcal{I}$

## Weak limits in $q\text{Cat}_2$

$$\begin{aligned}
 \mathcal{F}: q\text{Cat}_2(A, X^{\mathcal{P}}) &\cong h((X^{\mathcal{P}})^A) \cong h((X^A)^{\mathcal{P}}) \stackrel{?}{\cong} (h(X^A))^{\mathcal{P}} \\
 &\cong (q\text{Cat}_2(A, X))^{\mathcal{P}}
 \end{aligned}$$

$\uparrow$  cotensors commute with internal hom

It defines a notion of cotensoring by  $\mathcal{P}$

- We require  $h(X^{\mathcal{P}}) \cong (hX)^{\mathcal{P}} \longrightarrow \bullet$
- In  $q\text{Cat}_\infty$   $X^{\Delta^1} \in q\text{Cat}$ . Also,  $\mathcal{N}\mathcal{P} = \Delta^1$  and  $h\Delta^1 = \mathcal{P}$   
 $h(X^{\Delta^1}) \longrightarrow (hX)^{h\Delta^1} \cong (hX)^{\mathcal{P}}$  is not iso

• A weak cotensor with  $\mathcal{Z}$

Lemma (the universal property) The canonical comparison functor

$$h(X^{\Delta^1}) \rightarrow (hX)^{\mathcal{Z}}$$

is is surjective on objects

full

&

conservative

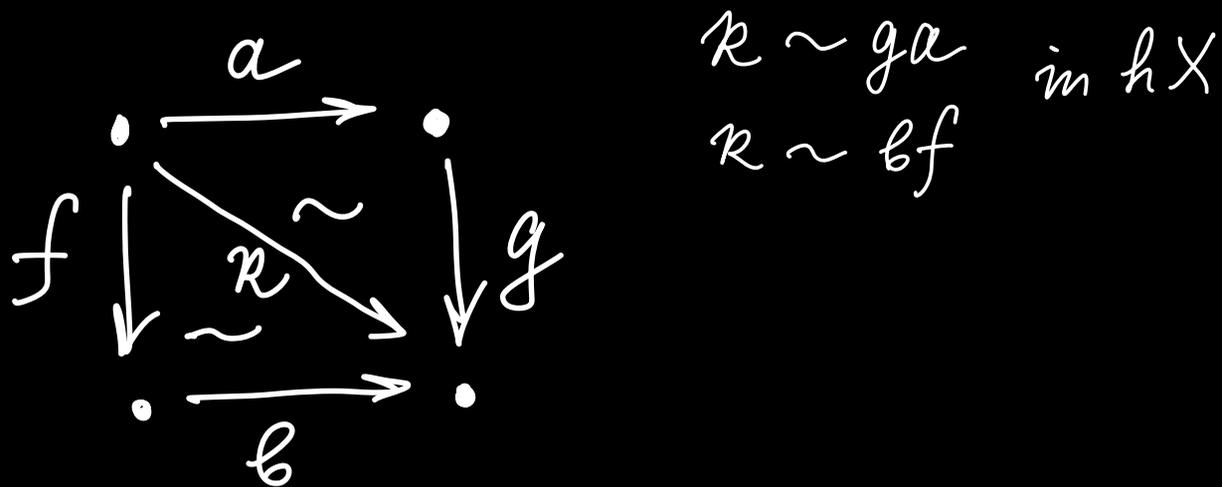
reflects iso

— it's smothering

▷ • Surjectivity: every arrow in  $hX$  is represented by a 1-simplex in  $X$

• Fullness: find a morphism from  $h(X^{\Delta^1})$  to

$$\begin{array}{ccc} \bullet & \xrightarrow{a} & \bullet \\ f \downarrow & & \downarrow g \\ \bullet & \xrightarrow{b} & \bullet \end{array}$$



We have

$$\Delta^1 \times \Delta^1 \longrightarrow X \iff \Delta^1 \longrightarrow X^{\Delta^1}$$

It represents the desired arrow in  $h(X^{\Delta^1})$

- Conservativity: omit it!

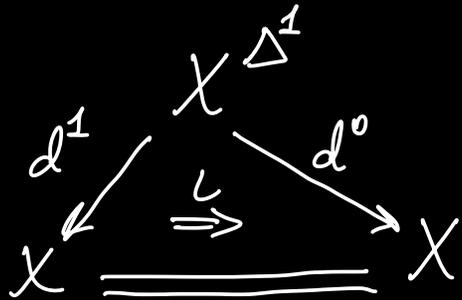


- $q\text{Cat}_2(A, X^{\Delta^1}) \cong h((X^{\Delta^1})^A)$

$$\cong h((X^A)^{\Delta^1}) \rightarrow (h(X^A))^{\mathcal{P}} = q\text{Cat}_2(A, X)^{\mathcal{P}}$$

natural in  $A$

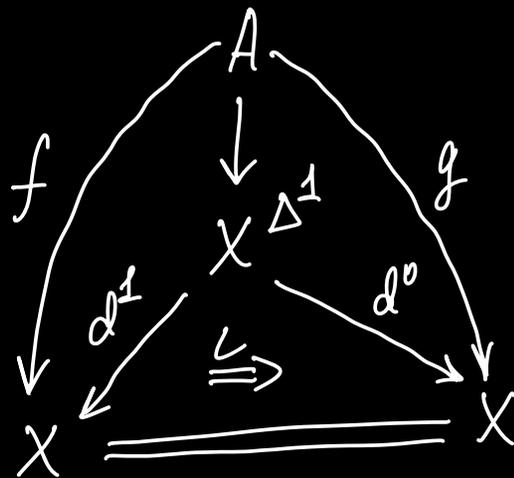
- Put  $A = X^{\Delta^1}$  and image of id will be:



In general case:

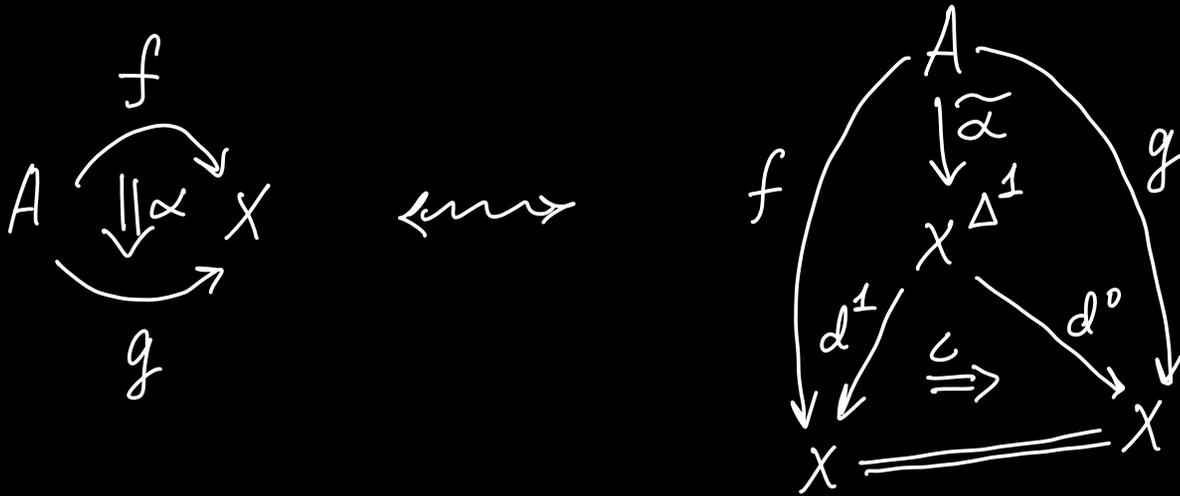
$$A \rightarrow X^{\Delta^1} \in \text{Ob}(\text{qCat}(A, X^{\Delta^1}))$$

a homotopy from  $f$  to  $g$



the composite 2-cell

Surjectivity of  $q\text{Cat}_2(A, X^{\Delta^1}) \rightarrow (q\text{Cat}_2(A, X))^2$  says:



$$l \cdot \tilde{\alpha} = \alpha$$

whiskering

- $\tilde{\alpha}$  is not unique:  $X^{\Delta^1}$  is only a weak cotensor by  $\mathcal{C}$
- The universal property defines the arrow quasi-categories up to equivalence

Lemma Let  $\mathcal{Z}$  be in  $\mathcal{Q}Cat$ , s.t.  $\exists$  natural transformation

$$h(\mathcal{Z}^A) \rightarrow (h(X^A))^{\mathbb{Z}}$$

is smothering (i.e., surjective on objects, full and conservative)

Then

$$\mathcal{Z} \cong X^{\Delta^1}$$

$\triangleright$  •  $\exists A = \mathcal{Z}$ , the image of  $1_{\mathcal{Z}} \in h(\mathcal{Z}^{\mathcal{Z}})$  is

$$\begin{array}{ccc} & \mathcal{Z} & \\ e^1 \swarrow & \mathcal{R} \Rightarrow & \searrow e^0 \\ X & \xlongequal{\quad} & X \end{array}$$

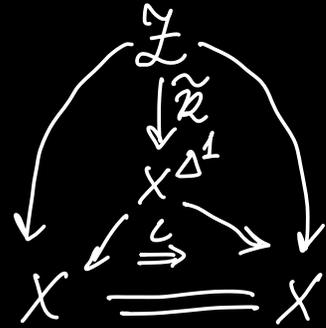
• Surjectivity implies:

$$A \begin{array}{c} \xrightarrow{f} \\ \Downarrow \alpha \\ \xrightarrow{g} \end{array} X \rightsquigarrow$$

$$\begin{array}{ccc} & A & \\ & \downarrow \alpha & \\ & \mathcal{Z} & \\ e^1 \swarrow & \mathcal{R} \Rightarrow & \searrow e^0 \\ X & \xlongequal{\quad} & X \end{array}$$

- Apply the weak univ. prop. of  $\mathcal{L}$  to  $k$ :

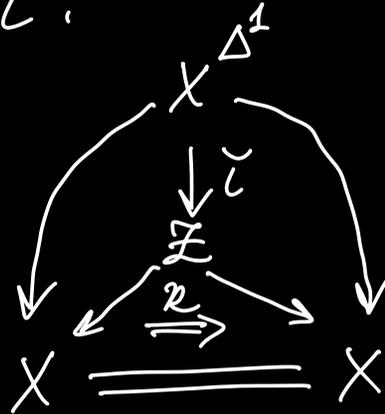
$$\tilde{\mathcal{R}}: \mathcal{Z} \rightarrow X^{\Delta^1}$$



$$\mathcal{L} \cdot \tilde{\mathcal{R}} = \mathcal{R}$$

- Apply the univ. prop. of  $\mathcal{R}$  to  $L$ :

$$\tilde{\mathcal{L}}: X^{\Delta^1} \rightarrow \mathcal{Z}$$



$$\mathcal{R} \cdot \tilde{\mathcal{L}} = L$$

$$X^{\Delta^1} \xrightarrow{\tilde{\mathcal{L}}} \mathcal{Z} \xrightarrow{\tilde{\mathcal{R}}} X^{\Delta^1}$$

gives a factorization  $L$  through itself

$$\mathcal{F}(\tilde{\mathcal{R}}\tilde{\mathcal{L}}) \cong \mathcal{F}(1_{X^{\Delta^1}}) \xRightarrow{\text{by conservativity \& fullness}} \tilde{\mathcal{R}}\tilde{\mathcal{L}} \cong 1_{X^{\Delta^1}} \text{ in } h((X^{\Delta^1})^{X^{\Delta^1}})$$

$$\mathcal{F}: h((X^{\Delta^1})^{X^{\Delta^1}}) \rightarrow h(X^{(X^{\Delta^1})})^{\mathcal{Z}}$$

• Similarly,

$$\tilde{\mathcal{R}} \cong \mathbb{1}_{\mathcal{Z}} \text{ in } h(\mathcal{Z}^{\mathcal{Z}})$$

• These iso's are represented by

$$\mathcal{Y} \rightarrow (X^{\Delta^1})^{X^{\Delta^1}} \quad \& \quad \mathcal{Y} \rightarrow \mathcal{Z}^{\mathcal{Z}}$$

where  $\mathcal{Y} = \mathcal{N}(\cdot \cong \cdot)$

and

Recall  $f: \Delta^1 \rightarrow X$  is an iso in a quasi-category  $\Leftrightarrow \exists$  an extension to  $\mathcal{Y}$   
to  $\mathcal{Y}$

• So, we have  $X^{\Delta^1} \times_{X^{\Delta^1}} X^{\Delta^1} \xrightarrow{\cong} X^{\Delta^1}$  by adjunction △

$$\mathcal{Z} \times_{\mathcal{Z}} \mathcal{Z} \xrightarrow{\cong} \mathcal{Z}$$

Remark It can be generalized to any category  
freely generated by a graph categories

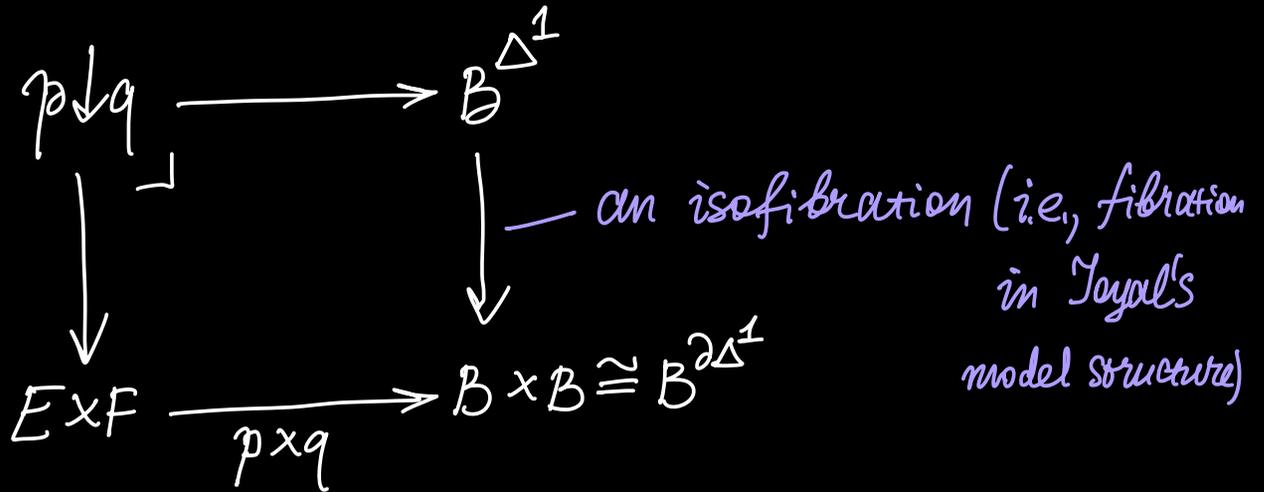
Lemma  $\forall$  diagram  $E \xrightarrow{p} B \xleftarrow{q} F$  a fibration in Joyal's model  
structure  
of cofibrant objects  
of quasi-categories with  $q$  an isofibration

$$h\left(\begin{array}{c} E \times F \\ B \end{array}\right) \longrightarrow \begin{array}{c} hE \times hF \\ hB \end{array}$$

is smothering functor.

## Comma quasi-categories

$E \xrightarrow{p} B \xleftarrow{q} F$  is a diagram of quasi-categories



Corollary The canonical functor

$$h(p \downarrow q) \rightarrow h(E \times F) \times_{h(B \times B)} h(B^{\Delta^1}) \rightarrow \underbrace{(hE \times hF) \times (hB)^2}_{hB \times hB} = h(p) \downarrow h(q)$$

is smothering

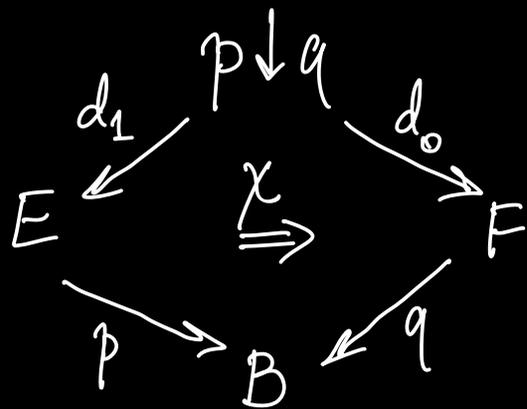
~~~~~  
 this is the usual comma category in Cat

- The functor

$$h((p \downarrow q)^A) \rightarrow h(E^A) \times_{h(B^A) \times h(B^A)} h(F^A) \times h(B^A)^2 = h(p^A) \downarrow h(q^A)$$

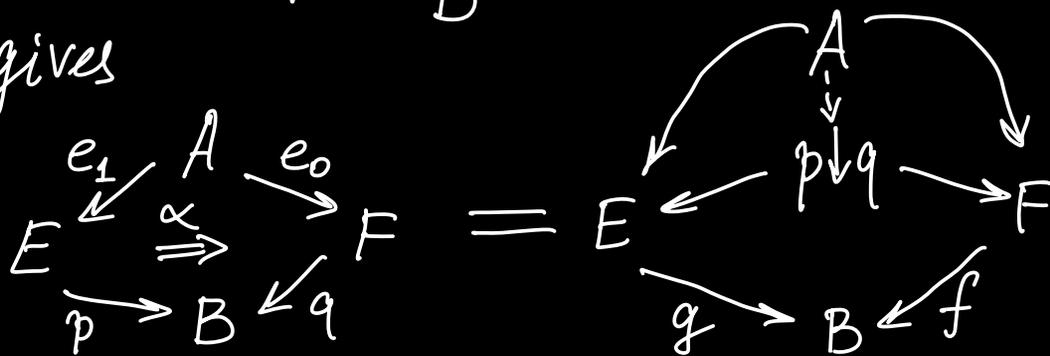
is also smothering

- The Weak universal property of  $p \downarrow q$



— the image of the identity at  $p \downarrow q$

Surjectivity gives



- By fullness & conservativity, one can derive that

$$\text{if } \gamma \cdot \left( f, g : A \rightrightarrows p \downarrow q \right) = \alpha$$

*whiskering* (with arrow pointing to the dot)

Then  $\exists$  an iso

$$A \begin{array}{c} \xrightarrow{f} \\ \Downarrow \cong \\ \xrightarrow{g} \end{array} p \downarrow q$$

represented by a map  $A \times \mathcal{I} \rightarrow p \downarrow q$

- $\Rightarrow p \downarrow q$  is unique up to equivalence

## The definition of $\infty$ -cosmos

Def. An  $\infty$ -cosmos is a simplicially enriched category  $K$ :

- objects by def. are  $\infty$ -categories
- hom's are quasi-categories
- there are a subcategory of isofibrations  $A \twoheadrightarrow B$

s.t. the following axioms hold:

(a) Completeness. -  $K$  possesses a terminal object  $1$ ;  
cotensors  $A^U$ ,  $U \in \text{Set}$ ;  
pullbacks of isofibrations along any functor

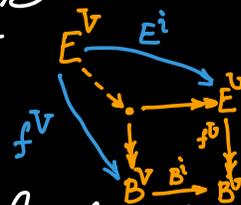
(b) Isofibrations. - The class of isofibrations  $\supset$  isomorphisms

$$!: A \overset{U}{\twoheadrightarrow} 1$$

- This class is stable under pullback along all functors

-  $\exists p: E \twoheadrightarrow B, i: U \hookrightarrow V \Rightarrow i \hat{\wedge} p: E^V \twoheadrightarrow E^U \times_{B^U} B^V$

the Leibnitz cotensor



-  $\forall X \in \text{Ob}(K), \forall p: E \twoheadrightarrow B \Rightarrow \text{fun}(X, p): \text{fun}(X, E) \twoheadrightarrow \text{fun}(X, B)$

(c) Equivalences.  $f: A \xrightarrow{\sim} B$  is an equivalence when

$$\text{fun}(X, f): \text{fun}(X, A) \twoheadrightarrow \text{fun}(X, B)$$

is an equivalence of quasi-categories  $\forall X \in K$

(d) Cofibrancy. All objects (say,  $A$ ) are cofibrant

$$\begin{array}{ccc} & \exists & E \\ & \nearrow & \downarrow \wr \\ A & \longrightarrow & B \end{array}$$

From the axioms above one can derive the stability of  $\text{Fib}^{tz}$

(d) Trivial fibrations. — They are defined to be in  $\text{WF} \cap \text{Fib}$

— They define a subcat  $\supset$  Isomorphisms

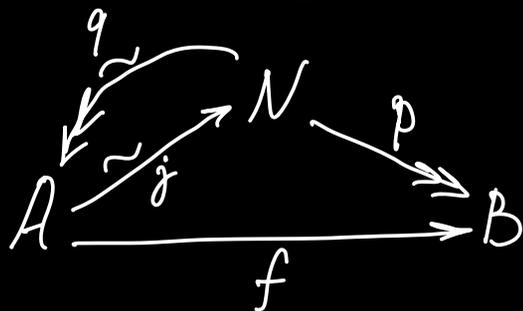
— They are stable under pullback along all functors

$$i \pitchfork p: E^V \longrightarrow E^U \times_{B^U} B^V \in \text{Fib}^{tz} \quad \begin{array}{l} p: E \twoheadrightarrow B \\ i: U \hookrightarrow V \end{array}$$

When  $p \in \text{Fib}^{tz}$  or  $i \in \text{Cofib}^{tz}$  in Joyal model structure on  $\text{sSet}$

Brown factorization lemma

(e) Factorization.  $\forall$  functor  $f: A \rightarrow B$   $f = p \circ j$



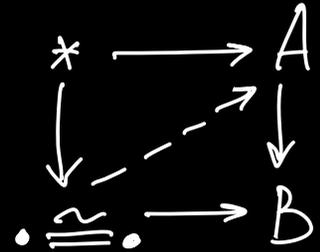
(f) Cartesian closure.  $- \times - : \mathcal{K} \times \mathcal{K} \rightarrow \mathcal{K}$  extends to

$$\text{fun}(A \times B, C) \cong \text{fun}(A, C^B) \cong \text{fun}(B, C^A)$$

## Examples

$M_{cf}$  enriched over the Joyal model structure on  $sSet$  defines an  $\infty$ -cosmos

- $Cat$  — the  $\infty$ -cosmos of small categories
- $Isfib$  — the usual isofibrations of cats
- $Equiv$  — the usual equivalences of cats
- $qCat$  — the  $\infty$ -cosmos of quasi-cat
- $CSS$  — the  $\infty$ -cosmos complete Segal spaces
- $Segal$  — the  $\infty$ -cosmos of Segal cats
- $sSet_+$  — the  $\infty$ -cosmos of naturally marked simplicial sets



•  $\mathbb{H}_n\text{-Sp}$  — the  $\infty$ -cosmos of  $\mathbb{H}_n$ -spaces, a simplicial presheaf model of  $(\infty, n)$ -categories

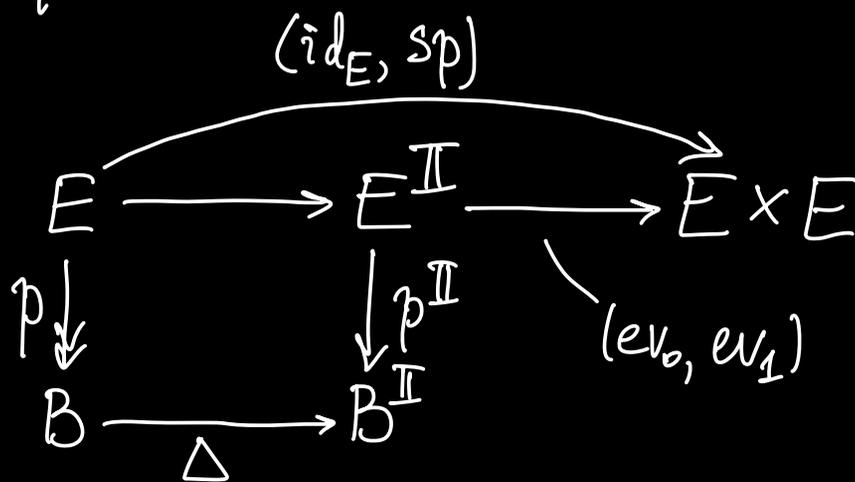
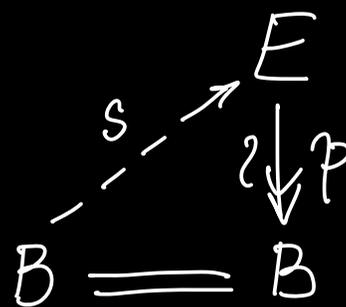
•  $\text{Rezk}_{\mathcal{M}}$  — the  $\infty$ -cosmos of Rezk objects in a nice model cat  $\mathcal{M}$

They are used to define iterated Segal spaces — another simplicial presheaf model of  $(\infty, n)$ -categories

Lemma. The equivalences in an  $\infty$ -cosmos are closed under retraces and satisfy 2-of-3 property

Lemma (trivial fibrations split)

The section defines a split fiber homotopy sequence

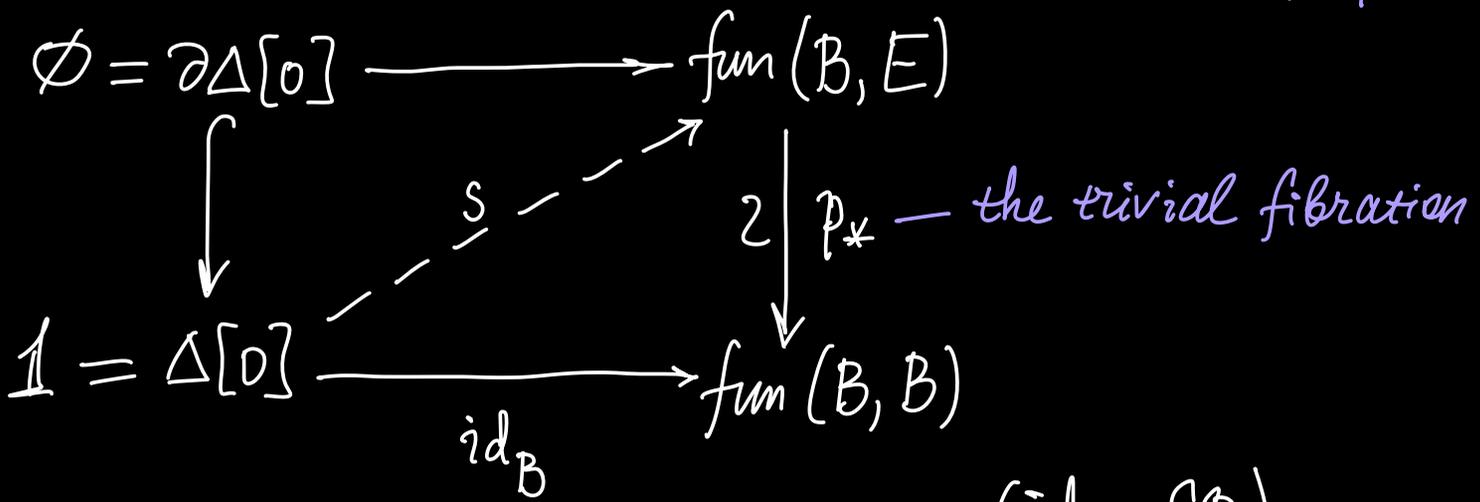


and conversely,  $\forall$  isofibration that defines a split fiber homotopy equivalence is a trivial fibration

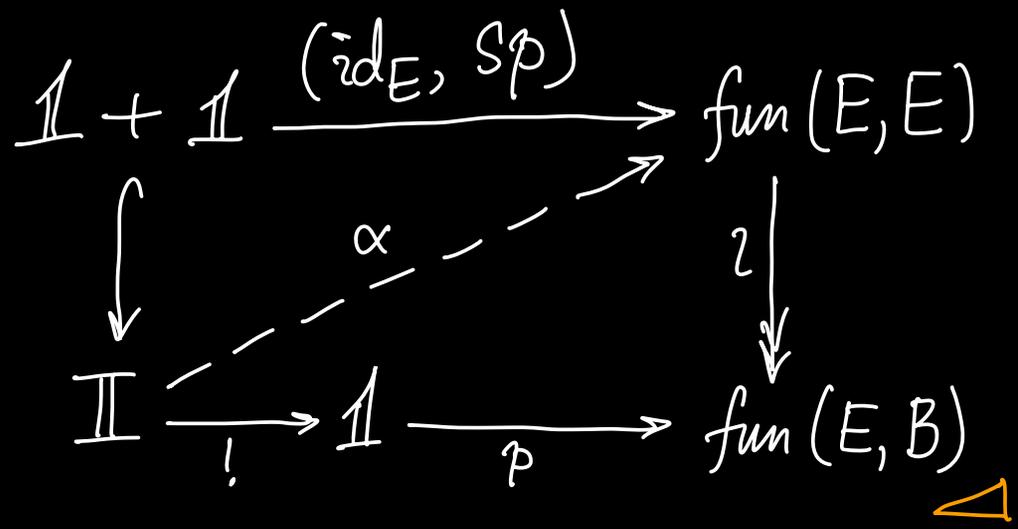
▷ If  $p: E \xrightarrow{\sim} B$  is trivial  $\Rightarrow p_*: \text{fun}(X, E) \xrightarrow{\sim} \text{fun}(X, B)$

$\forall \infty\text{-cat } X$

by the stability properties of  $\text{Fib}^{\text{tr}}$



After that we can solve the problem



## Recall: Formal category theory in a 2-category

- Objects are called  $\infty$ -categories
- 1-cells  $f: A \rightarrow B$  is said to be  $(\infty)$ -functors
- 2-cells  $A \begin{array}{c} \xrightarrow{f} \\ \Downarrow \alpha \\ \xrightarrow{g} \end{array} B$  -  $(\infty)$ -natural transformations

Def. An adjunction between  $\infty$ -categories:

- $f: B \rightarrow A, u: A \rightarrow B$
- $\eta: id_B \Rightarrow uf$  &  $\epsilon: fu \Rightarrow id_A$

$$\begin{array}{c}
 B \xrightarrow{u} B \\
 \uparrow f \quad \downarrow u \\
 A \xrightarrow{u} A
 \end{array}
 = u \left( \begin{array}{c} B \\ \xrightarrow{id} B \\ \Downarrow idu \\ A \end{array} \right) u
 \quad
 \begin{array}{c}
 B \xrightarrow{u} B \\
 \downarrow f \quad \uparrow u \\
 A \xrightarrow{u} A
 \end{array}
 = f \left( \begin{array}{c} B \\ \xrightarrow{id} B \\ \Downarrow idf \\ A \end{array} \right) f$$

Prop. Adjunctions compose:

$$C \begin{array}{c} \xleftarrow{f'} \\ \perp \\ \xrightarrow{u'} \end{array} B \begin{array}{c} \xleftarrow{f} \\ \perp \\ \xrightarrow{u} \end{array} A \rightsquigarrow C \begin{array}{c} \xleftarrow{ff'} \\ \perp \\ \xrightarrow{u'u} \end{array} A$$

Def. An equivalence between  $\infty$ -categories consists of:

- a pair of  $\infty$ -categories  $A$  and  $B$
- $f: A \rightarrow B$  &  $g: B \rightarrow A$  — functors

• 
$$A \begin{array}{c} \xrightarrow{f} \\ \cong \\ \downarrow \alpha \\ \xrightarrow{g} \end{array} A \quad \& \quad B \begin{array}{c} \xrightarrow{fg} \\ \cong \\ \downarrow \beta \\ \xrightarrow{gf} \end{array} B$$

# The homotopy 2-category of an $\infty$ -cosmos

Def. This is a strict 2-cat  $K_2$  or  $hK$  so that

- $Ob(K_2) = Ob(K) = \infty\text{-categories}$
  - 1-cells  $f: A \rightarrow B$  of  $K_2 \iff$  the vertices  $f \in \text{fun}(A, B)$
  - $\infty$ -functors  $f$
  - 2 cells  $A \begin{array}{c} \xrightarrow{f} \\ \Downarrow \alpha \\ \xrightarrow{g} \end{array} B$  in  $K_2 \iff$  homotopy classes of 1-simplices
- $\alpha: f \longrightarrow g \in \text{fun}(A, B)$
- $K_2 := h_* K$
- $2\text{-Cat} \xleftarrow{h_*} \text{Set-Cat}$
- The cat  $qCat_2$  was first introduced by Joyal

Or, just  $h\mathcal{K}$ :

$$\bullet \text{ob}(h\mathcal{K}) := \text{ob}(\mathcal{K})$$

$$\bullet \text{hom}(A, B) := h(\text{fun}(A, B))$$

homotopy cat of quasi-cat

Def. The underlying cat of 2-cat — simply forgetting  
its 2-cells

$$u(\mathcal{K}) = u(h\mathcal{K})$$

Denote by  $\mathcal{Q} := \mathcal{N}(\cdot \longrightarrow \cdot) = \Delta^1$

$$\begin{array}{c}
 f \\
 \curvearrowright \\
 A \quad \Downarrow \alpha \quad B \\
 \curvearrowleft \\
 g
 \end{array}
 \iff \alpha : \mathcal{Q} \longrightarrow \text{fun}(A, B)$$

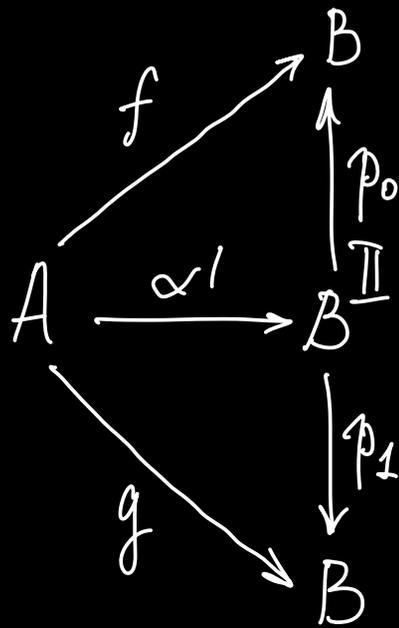
Transpose:

$$\begin{array}{ccc}
 & & B \\
 & \nearrow f & \\
 A & \xrightarrow{\alpha} & B^{\mathcal{Q}} \\
 & \searrow g & \\
 & & B \\
 & & \downarrow p_1 \\
 & & B \\
 & & \uparrow p_0
 \end{array}$$

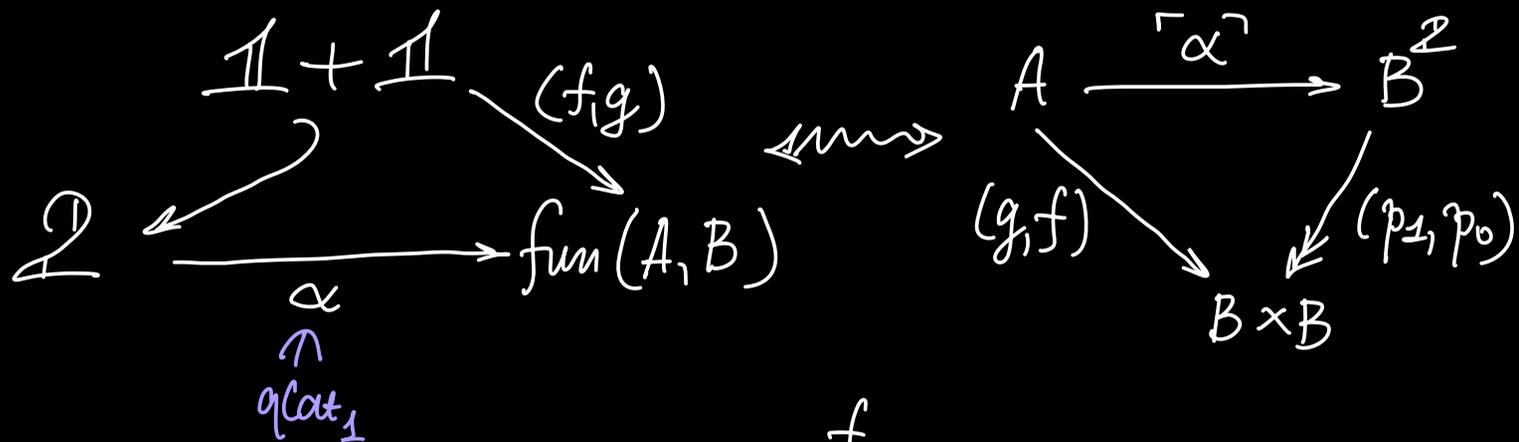
$$\begin{array}{c}
 f \\
 \curvearrowright \\
 A \xrightarrow{\quad} B \\
 \Downarrow \alpha \\
 \curvearrowleft \\
 g
 \end{array}
 \text{ is an iso in } \mathcal{K}_2 \iff \text{ho}(\alpha): \text{ho}\mathcal{Z} \rightarrow \text{ho fun}(A, B) \text{ is iso}$$

$$\iff \alpha: \mathcal{Z} \rightarrow \text{fun}(A, B) \text{ extends to } \alpha': \mathbb{I} \rightarrow \text{fun}(A, B)$$

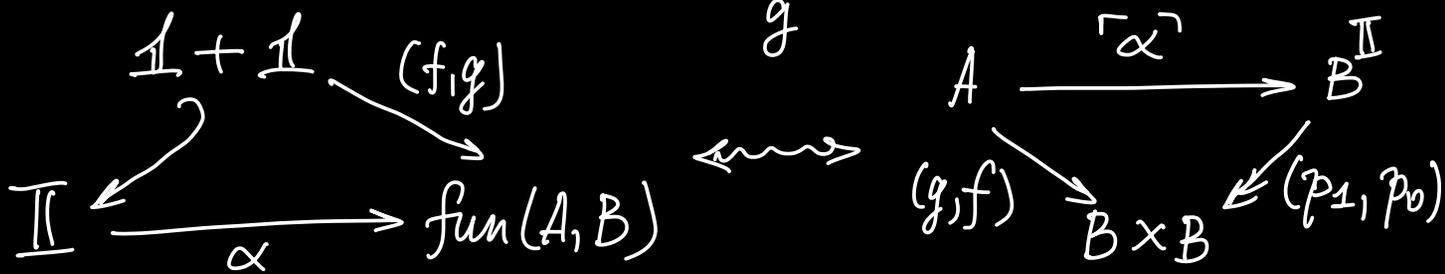
$$\mathcal{N}(\cdot \cong \cdot)$$



Lemma. (i)  $\forall$  2-cell  $A \begin{matrix} \xrightarrow{f} \\ \Downarrow \alpha \\ \xrightarrow{g} \end{matrix} B$  in  $\mathcal{K}_2$   
 is repr. by a map in  $qcat$  the homotopy 2-cat of an  $\infty$ -cosmos



(ii)  $\forall$  invertible 2-cell  $A \begin{matrix} \xrightarrow{f} \\ \Downarrow \cong \alpha \\ \xrightarrow{g} \end{matrix} B$  in  $\mathcal{K}_2$  — " —



## Functors between $h\mathcal{K}$

Lemma.  $\forall$  simplicial functor  $F: \mathcal{K} \rightarrow \mathcal{L}$   <sup>$\infty$ -cosmoi</sup> induces  
 a  $\mathcal{L}$ -functor  $F: \mathcal{K}_2 \rightarrow \mathcal{L}_2$

▷ On objects and 1-cells — it's okay

$$\begin{array}{ccc}
 \text{fun}(A, B) & \xrightarrow{F} & \text{fun}(FA, FB) \\
 \begin{array}{ccc}
 \xrightarrow{f} & & \\
 \Downarrow \alpha & & \\
 \xrightarrow{g} & & 
 \end{array} & \longmapsto & \begin{array}{ccc}
 \xrightarrow{Ff} & & \\
 \Downarrow F\alpha & & \\
 \xrightarrow{Fg} & & 
 \end{array}
 \end{array}$$

$F$  is a morphism of simplicial sets

So, homotopic 1-cells map to homotopic ones



## Cartesian closure and products

- Prop.
- (i)  $\mathcal{K}_2$  has 2-categorical products
  - (ii)  $\mathcal{K}_2$  is cartesian closed as 2-cat

Theorem (Equivalences of  $\infty$ -categories are 2-cat equiv.)

In any  $\infty$ -cosmos  $\mathcal{K}$  the F.A.E.:

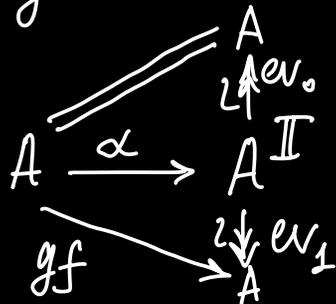
(i)  $\forall X \in \mathcal{K}$

$$f_x: \text{fun}(X, A) \xrightarrow{\sim} \text{fun}(X, B)$$

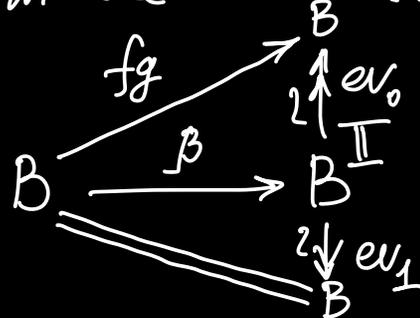
— equiv. of quasi-cat

(ii)  $\exists g: B \rightarrow A$  &  $\alpha: \text{id}_A \cong gf$ ,  $\beta: fg \cong \text{id}_B$  in the homotopy 2-cat

(iii)  $\exists g: B \rightarrow A$  and maps in the  $\infty$ -cosmos  $\mathcal{K}$



&



Thank you!